

**PROGRAM LISTING ORGANIZATION AND USAGE**  
**COMMON SYSTEMS SOFTWARE DESCRIPTION**  
**3A PROCESSOR**

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**1. GENERAL**

**1.01** This section describes the format and types of information contained in program listings (PRs) for the 3A Processor. It also provides an aid for using the program listings. In addition to this document, the following sections will aid in understanding and using the program listings:

SECTION	TITLE
232-160-100	Introduction to No. 2B Electronic Switching System, Software General Description, Generic EF-2
232-164-110	Program Listing Standards, Support Software Description, No. 2B Electronic Switching System
232-164-115	Basic Program Instruction Set Description, No. 2B Electronic Switching System
254-340-082	System Utilities, Software Subsystem Description, 3A Processor Extended Operating System
254-340-100	Introduction to 3A Language, 3A Processor Common Systems
254-340-102	Basic and Extended 3A Instruction Set, 3A Processor Common Systems

**1.02** Whenever this section is reissued, the reason for reissue will be listed in this paragraph.

**1.03** A list of abbreviations and acronyms used in this section is provided in Table A.

**2. PROGRAM LISTING DESCRIPTION**

**PROGRAM DESIGN**

**2.01** Modular program design is used for the 3A Processor programs. The structure of a program is modular when the different tasks, subtasks, or functions are divided logically and packaged into relatively small program units which perform well-defined functions. Large programs

[assembly unit or program identification (PIDENT), depending on the system] are divided into several small sections or program units (PUs) for ease of design, assembly, administration, maintenance, and readability. The assembly unit is the largest physical portion of a modular program, and is compiled or assembled as one entity. The assembly unit may or may not contain functionally related program units. A program unit is a subportion of a modular program and is a collection of code within an assembly unit which performs a well-defined function and is given a name by which it may be called. The programming techniques used also make use of pseudo-operations, macros, and subroutines in designing programs.

**2.02** A pseudo-operation is an operation which may be used in an assembly statement to control assembler activities but does not result in executable code. These are instructions to the assembly program which control the structuring of information found in the listing. Pseudo-operations are an aid in writing and checking the program, in making storage-address assignments, and in defining data. PUBLIC, EXTERN, EQU, and EQU(A) are examples of pseudo-operations.

**2.03** A macro is a sequence of operations called by an abbreviated label or name. The macro can generate sequences of code depending on the parameters supplied in the macro call. The coded lines generated by the macro are easily identifiable in a listing by the macro level number which is surrounded by dashes, eg, -001- (macro level 1) or -002- (macro level 2). Macros may be contained within the PIDENT or may be contained in an external PIDENT referenced by the pseudo-operation GETLIB.

**2.04** A subroutine is a block of code which provides a facility for maintaining a single copy of a given routine which can be used in a repetitive manner by other programs located in memory. Control of processing is passed from the calling program/subroutine to the called subroutine until the function is completed. Subsequently, control is returned to a location specified by the calling program. Subroutines may be either common or private subroutines. A common subroutine is one which has been declared PUBLIC; it can be used by a program specifying its use by an EXTERN pseudo-operation. It thus may be accessed by many programs. A private subroutine generally has application only in the PIDENT in

which it is listed and is not declared as a common subroutine by a PUBLIC pseudo-operation; thus, is not accessible by other programs.

**2.05** In using common subroutines, interface requirements must be satisfied. Verification of subroutine interfaces during assembly can be accomplished with macros. Various macros are provided by different systems to implement calling subroutines and returning from subroutines. These macros may have slightly different forms in the various applications; however, in most cases, the forms will incorporate the words CALL and RETURN either alone or with additional suffixes or prefixes added, eg, RETURNOS, S\_CALL, SRETURN, etc.

#### Conversion Between Binary, Octal, and Hexadecimal Numbers

**2.06** All digital computers use the binary system (two-state elements) internally to perform operations. A binary number is expressed in zeros and ones, such as: 1101. A binary digit is a number that can have one of two unique values; 0 or 1 and no other value. Binary digits can be used to represent numbers of any magnitude just as a string of decimal digits can be used to represent numbers greater than nine (9). A decimal digit can have one of ten unique values; 0 through 9. Decimal numbers cannot be easily handled by a computer, therefore decimal numbers are ultimately converted to binary numbers for use in a computer. Decimal numbers are called base ten numbers since there are ten unique values (0 through 9), and binary numbers are called base two numbers since there are only two unique binary values. Because binary numbers tend to be very long, binary numbers are often converted to another base number system. Conversion is accomplished by grouping the binary number into sets or groups of three or four binary digits and then converting to the other system. The numbering systems most often used are the base eight (three binary digits are used to represent eight unique values, 0 through 7), or base sixteen (four binary digits are used to represent sixteen unique values, 0 through 9 plus the letters A through F following after the digit 9). The octal or hexadecimal (base eight or base sixteen, respectively) numbers are often used rather than the base two (binary) numbers.

**2.07** The 3A Processor uses either a 16-bit or 24-bit binary word, depending on the application. Therefore, a shorthand method called

the octal number system is used (the hexadecimal number system is also used) to express fields in the listings. When each binary word is divided into 3-bit groups starting at the right, each group can be represented by one of the eight octal numerals 0, 1, 2, 3, 4, 5, 6, and 7.

**2.08** Conversion between binary and octal numbers (base two to base eight) can be accomplished easily. The binary number is divided into sets or groups of three binary digits (beginning with the rightmost binary digit). Each position in the set of three represents a power of the numeral 2 whenever that position is occupied by a binary 1, and if a binary 0 occupies a position, it represents a zero. Hence the position of a binary 1 in the set of three binary digits represents a power of two that corresponds to the position of the binary digit 1. The rightmost position represents the zero power, the middle position represents the first power, and the leftmost position represents the second power of the numeral 2. The octal number is the sum of the powers of 2 denoted by the binary ones that are set in each group of three binary numbers when the word bits are divided into groups of three, beginning with the rightmost bit. For example, binary 001 represents the zero power of 2 which is octal number 1, binary 100 represents the second power of 2 which is octal number 4, and binary 111 represents the second, first, and zero powers of 2 which is octal number 7 (the sum of those 3 powers of 2, eg, 4 plus 2 plus 1). Binary - octal conversion is as follows:

Binary:

000 001 010 011 100 101 110 111

Octal:

0 1 2 3 4 5 6 7

**2.09** Examples of conversions are given below:

Binary to octal:

Binary:

1 110 010 101 001 011

Octal:

1 6 2 5 1 3

Octal to binary:

Octal:

0 2 6 4 1 5

Binary:

0 010 110 100 001 101

**2.10** Conversion between binary and hexadecimal (base sixteen) numbers is somewhat similar to the binary to octal conversion except; (1) the binary number is divided into sets of four binary digits rather than into sets of three binary digits, and (2) there are sixteen unique digits rather than eight as in the octal system. The position of a binary 1 in the set of four binary digits represents the power of two as before, except the highest power of two that can be represented by a binary 1 is now the third power of two rather than the second power. Table B provides a summary of the number system (binary, octal, hexadecimal, and decimal equivalents). The following is an example of a binary to hexadecimal conversion.

Binary:

0000 1100 0001 0010

Hexadecimal:

0 C 1 2

**2.11** A symbolic instruction, its octal representation (as it would appear in the listing), and the binary or hexadecimal equivalent are represented typically as follows:

Symbolic instruction:

LR R1 R2

Octal representation:

006022

Binary equivalent:

0000110000010010

Hexadecimal equivalent:

0C12

### 3A PROCESSOR PROGRAM LISTINGS

**2.12** A program listing (PR) is a hard copy (printout) of a program and is produced for each segment of code assembled as a unit by the assembler program (this unit may also be referenced as a PIDENT). Each listing is identified by a name consisting of a maximum of six characters and a PR number (not necessarily unique). The name is normally an acronym which is functionally relevant to the listing. The PR number provides identification for the listing within a particular software system (series of programs). The program listing normally contains a prologue and comments which explain the function and logic of a particular listing in addition to recording the actual coded instructions which perform the functions.

**2.13** A complete listing contains the following types of information. There is no absolute order in which the various types of information must appear; however, there is a general structure which can be followed. The various types of information found in the program listings are:

- Prologues and comments
- Control statements
- PUBLIC and EXTERN definitions
- Macro definitions
- Symbol definitions
- Data definitions
- Program section
- Symbol reference tables.

#### A. Prologues and Comments

**2.14** Prologues and comments are placed in a listing to clarify and explain the function of the program instructions. There is no standard prologue format; thus the prologue may appear in various formats and contain a variety of information. Prologues explain why the program instructions are necessary, what function is accomplished, why it is done, and when it is done in relation to the other program instructions or functions. The prologue may be in table form with the what, why, when, etc, listed explicitly in tabular format,

or the prologue may be a straight-forward narrative. Prologues are normally found at the beginning of a listing, but may be found throughout a listing. Comments are found throughout a listing and are used to explain one or more lines of code. Examples of various forms of prologues are shown in Fig. 1. Prologues and comments may be divided into four general categories:

- (a) An assembly unit prologue which, when appropriate, serves to define the purpose and function of the assembly unit (PIDENT).
- (b) A program unit prologue explains the function of well-defined functional program units within the assembly unit in greater detail than does the assembly unit prologue. The program unit prologues may be in various formats. Example prologues are shown in Fig. 2.
- (c) A comment is used to explain a line (or lines) of code. A comment is denoted by a sharp sign (#) which precedes the comment. Comments may apply to several lines (a block) of code. A block of code is a small section of instructions identifiable as performing a subfunction or operation and may be given a name and considered as an entity. An example of a comment on a block of code is shown in Fig. 3.
- (d) Comments may also be used on individual coded lines to explain the action of the code, data symbols, etc. An example of line comments is shown in Fig. 3.

## B. Control Statements

2.15 Control statements are used to direct the activity of the assembler, identify the input, and meet certain assembler requirements such as record keeping, etc. Certain mandatory pseudo-operations are required to satisfy both the assembler and the loader. Various control statements (Fig. 4) normally appear at the beginning of the listing and establish the identity and version of the PIDENT. This is accomplished by use of the NAME pseudo-operation and the VERSION macro. The various libraries and macros which are required for use by the PIDENT are identified by various forms of the GETxxx pseudo-operation, eg, GETLIB, GETMAC, etc. The program listing may be one for an ESS Programming Language (EPL) program

or one for a program which is written in assembler code. Control statements differ in the two cases:

- (a) For an EPL program, the first EPL statement will normally appear with the other control statements. This statement, BEGIN\_EPL, is used to initialize the EPL compiler and establish the various EPL options. Associated with the BEGIN\_EPL statement is an END\_EPL statement used to inform the compiler that no more EPL statements are to follow.
- (b) Programs written in assembly code need not have the BEGIN\_EPL statement.
- (c) The END\_EPL statement informs the compiler that no more EPL statements are to follow, ie, it is the limit of code to be assembled.

A detailed explanation of the control statements used in both types of listings is provided in Section 254-340-100.

## C. PUBLIC, EXTERN, and DIRECT Definitions

2.16 The declaration of PUBLIC and EXTERN definitions (Figs. 5 and 6) normally appears at the beginning of the listing but is not restricted to this location and may appear elsewhere in the listing. The PUBLIC pseudo-operation is used to declare those symbols (such as subroutine names, buffer addresses, and entry points) in a program that can be referenced and used by other programs. Examples of PUBLIC statements are shown in Fig. 5. Comments may be included to explain the meaning of the statements.

2.17 The EXTERN pseudo-operation is used in a program to declare a symbol which has been defined and declared PUBLIC in another program but must be referenced by this program. Examples of EXTERN statements are shown in Fig. 6. Comments may be included with a declaration to explain the meaning of the symbol. As with the PUBLIC statement, the location of the EXTERN statement is not restricted to the beginning of the listing.

2.18 The DIRECT pseudo-operation is used to declare a symbol which has been defined in another program but must be referenced by this program. The DIRECT pseudo-operation is defined in the program listed in the location symbol (Fig. 6) field of the DIRECT statement (similar to the

EXTERN statements). Comments may be included with this statement to explain the meaning of the symbol.

#### D. Macro Definitions

**2.19** When macros are defined in a program, the macro definition section may follow the PUBLIC and EXTERN definitions; however, the location is not restricted to the beginning of the listing. The macros that appear in a program listing are those used and/or controlled by that program to serve some specialized purpose. However, when a macro is used by more than one program, its definition is usually placed in a public macro library accessible by all programs to avoid duplication in coding or changing the macro. Each macro definition has the following elements:

- (a) Macro name (MACRO)
- (b) Parameters
- (c) Dummy variables
- (d) Functions and formula
- (e) Macro end (MEND).

Comments, preceded by a sharp (#) sign, may be included to explain the function and purpose of the macro. An example of a macro is shown in Fig. 7.

#### E. Symbol Definition

**2.20** In most cases pseudo-operations are used to define symbols when the symbol is placed in the location field. There is a collection of pseudo-operations (EQU, EQUR, EQURA, and SET) which are used solely for defining symbols. These may be used to equate two symbols, define program parameters, and assign symbolic values to general registers. This set of pseudo-operations normally appears after the PUBLIC and EXTERN definitions in the listing.

**2.21** The EQU pseudo-operation permanently defines the symbol in the location field. Symbol is the name of the symbol to be defined and value is any expression whose value is assigned to symbol. Value may be a decimal, binary, octal, or hexadecimal number, or another symbol. This

operation is used primarily for constants or program parameters.

LOCATION	OP	OPERAND
symbol	EQU	value

where value is the expression in the operand field. The SET pseudo-operation is identical to EQU except that it may be used to redefine a symbol without causing a multiple definition.

**2.22** The 16 general registers are identified by numbers 0 through 15. However, the registers can be named symbolically and should be when this clarifies the use of the registers. The EQUR pseudo-operation is used to permanently define the symbols used for registers. Symbol is the name of the symbol to be defined by a decimal arithmetic expression value in the operand field and may have a numeric value of 0 through 15.

LOCATION	OP	OPERAND
symbol	EQUR	value

where value is any decimal arithmetic expression whose value is 0 through 15 inclusive.

**2.23** Address register pairs RA0 and RA1 are defined symbolically by the EQURA pseudo-operation. Symbol in the location field is the name of the symbol which is defined by value in the operand field. The arithmetic expression value must be either 12 or 14 which will define the register pair 12 (register pair 12 and 13) or the register pair 14 (register pair 14 and 15) as the address register pair represented by symbol.

LOCATION	OP	OPERAND
symbol	EQURA	value

where value is any decimal arithmetic expression whose value is 12 or 14. Figure 8 provides examples of symbol definitions for registers.

**2.24** The TEXT pseudo-operation is used to define a symbol to be equivalent to a string of text.

LOCATION	OP	OPERAND
symbol	TEST	"Text string"

where symbol is the TEXT symbol being defined and "Text string" are the characters to be used as the value or contents of the text symbol. The first and last quotes on the "Text string" serve as delimiters for the text string. The string itself may contain quotes.

**2.25** The SYN (synonym) pseudo-operation is used to define a symbol (symbol1) as a synonym of another symbol (symbol2). The relationship is a lasting or permanent one even if the first symbol is redefined (using the SET pseudo-operation).

LOCATION	OP	OPERAND
symbol1	SYN	symbol2

where symbol1 is the name of the symbol being defined, and symbol2 is the name of the symbol synonymous with symbol 1.

**2.26** The ARRAY pseudo-operation is used to define an array of arithmetic values and to assign an initial value to the array elements. The individual elements of an array may be assigned values by the SET or the EQU pseudo-operations. An element of an array may be used anywhere a symbol may be used in an expression. The total size of an array should not exceed approximately 850 elements, and the limit of the array dimensions is 10.

LOCATION	OP	OPERAND
ARRAY	name	([a1:] b1, ..., [an:] bn) [=init], ..

where name is the name of the array, a1 is any decimal expression for the lower bound of the ith dimension, b1 is any decimal expression for the upper bound of the ith dimension, and init is the value assigned to each element of the array. The elements are referenced thus: name (ael,...,aen) where ael is an arithmetic expression.

## F. Data Definition Statements

**2.27** The data definition statements (Fig. 9) are listed next in the program listing, and the statements are normally structured using pseudo-operations and system macros. These statements provide symbolic definitions and layouts of the primary data elements (data blocks and words) used in the program listing. A detailed layout of a memory table (Fig. 10) will include the length (number of words) of the data block and the width (number of bits) of partitions within the block. A partition may be a word, a number of consecutive words, a bit, or a number of consecutive bits within the block that are assigned meaningful names. The named partitions are handled as separate entities by the program. As a further aid to understanding the layout, a picture of the data block may follow immediately after the data definitions. The most common statements used for data definitions are the system macros LAYOUT or LOUT, ITEM, LOSKIP, and LOEND.

- LAYOUT or LOUT is used to characterize and name a data block by giving a name to the block and defining the number of consecutive words assigned to the block. A block may consist of one or more words.
- ITEM is used to characterize and name the partitions (words or bits within a word) within the data block defined by LAYOUT. A partition within a data block may be a word (16 bits or 24 bits per word depending upon the word length capability of the particular processor), a group of bits, or single bits within a word which are assigned meaningful names. The naming of the words

or bits proceeds from the top of the item list to the bottom, and the naming of the bits proceeds from the rightmost to the leftmost bit within a 16-bit word. Thus, ITEM defines the width of the partitions of the data block proceeding from the first word location of the block to the last word location in an orderly manner.

- LOSKIP provides a means of automatically skipping words or bits in the data block (defined by LAYOUT) without counting the words or bits. The number of words to skip in the current block is defined or the number of bits to skip in the current word is defined.
- LOEND is used to indicate the end of the layout of the data block defined by LAYOUT or LOUT.

**2.28** There is also a data definition pseudo-operation which is commonly used for addresses, ie, ADDR.

LOCATION	OP	OPERAND
symbol	ADDR	value

The ADDR pseudo-operation can be used to generate a 20-bit address in a double word. The symbol can be defined with a value equal to the address word. Value can be expressed either as a relocatable address or as a nonrelocatable decimal expression. A special use of ADDR is as the operand in an ITEM macro statement, where it will automatically allocate the 20 bits of address. The subsequent ITEM statements should then allocate the remaining 12 bits of the first word of the address pair.

**2.29** Comments are often included in the data definition section to explain the use of the various data definitions. A comment normally follows the definition and is preceded by a sharp sign (#).

**G. Program Section**

**CSECT and DATASECT**

**2.30** The program section contains the program instructions and comments describing the

functions performed by the instructions. The program section may contain subsections called CSECT and DATASECT (Fig. 11). A CSECT is a pseudo-operation used to specify the beginning of a sequence of instructions which can be relocated as a block of instructions. A CSECT consists of a block of instructions that perform some well-defined function, identifies the beginning of an identifiable control unit, and assigns a location to the identity. CSECT is a mandatory pseudo-operation that must appear at the beginning of each control section in a program listing. Associated with the last CSECT is another mandatory pseudo-operation called END, which must appear at the end of the relocatable block comprising the assembly unit. The assembly unit may contain a number of program units (PUs) that perform some well-defined subfunctions of the assembly unit function. A DATASECT is a data CSECT used to declare areas of memory for the user program data. The DATASECT is given a label and appears at the beginning of the assembly unit.

**2.31** A program unit prologue (Fig. 1) is normally found at the beginning of each program unit and will contain an explanation of the program unit. The format of the prologue is not standard but will generally be in one of the formats described previously, either the narrative or the tabular format.

**2.32** An entry point is a labeled location at which a CSECT or a block of code may be entered and to which control may be transferred. An exit point is a location at which a block of code or CSECT terminates and gives control to some other location. Information about entry and exit conditions helps in the understanding of interrelationships between program units. Thus, comments may precede or accompany entry or exit points.

**2.33** Paragraphs 2.34 through 2.52 refer to the examples of program listing pages within program sections. The title of the program unit is given at the top left corner of each page with a subtitle below it, if applicable. The assembly unit or PIDENT name may be given at the top right corner of each page. Across the bottom of the page is information such as the PIDENT name, time and date of the listing, the PR number, issue number, and page number.

**2.34** Each program unit contains program instructions which instruct the 3A Processor to perform specific tasks. The program instructions are symbolic

instructions which will be translated into binary machine instructions. Refer to Section 254-340-102 for explanations of the instructions.

**2.35** The format of program instruction or line of code contains the following fields:

- Address field
- Encoded instruction field
- Branch address field
- Editor line number field or macro or EPL expansion level field
- Page line number field
- Location symbol field
- Symbolic operation code field
- Variable or operand field
- Comment field.

The fields are discussed in order of their appearance (left to right) on the program listing and are illustrated in Fig. 12.

#### Address Field

**2.36** The address field (Fig. 12) is reserved for the relative (CSECT) address (memory location) of the program instruction. The address field contains the memory address (expressed in octal or hexadecimal numbers) of that particular line of code. When this field is blank, the line has been used for some purpose which does not require a memory location, eg, macro calls, comments, etc. A double word instruction will cause the address field to be incremented by two memory locations. A double word instruction at location 00053732 will cause the next printed location to be 00053734.

#### Encoded Instruction Field

**2.37** The encoded instruction field entries are the octal or hexadecimal equivalent of the binary machine code, translated from a symbolic instruction. Only executable code entries appear in this field. No entries are present for a macro call, comments, or EPL statements, but are present for the macro or EPL expansions when they are

listed. When expansions are suppressed in an EPL program, this field is blank (see Fig. 13) for all instructions to follow until the next EPL options statement is encountered that turns the EXPAND option on. The expansion is suppressed or not suppressed by setting the EXPAND option in the BEGIN\_EPL statement to OFF or to ON respectively. For example, the statement

```
BEGIN_EPL EXPAND=OFF
```

suppresses the printout of the encoded instruction, whereas the statement

```
BEGIN_EPL EXPAND=ON
```

does not suppress the printout of the encoded instruction. There are variations in the encoded instruction field. These statements apply to the whole assembly unit. The EPL options statement EPLOPTIONS EXPAND = ON and EPLOPTIONS EXPAND = OFF are similar to the above but apply within an assembly unit rather than to the whole assembly unit. The field is different for the assembly code instructions, the 3A Processor instructions, and the 2B ESS instructions, (see Figs. 12 and 14).

**2.38** For the assembly code (Fig. 12) the first digit in the encoded instruction field represents the branch-allowed (BA) bit. (The BA bit is equivalent to the TA bit used with the 2B ESS instructions.) When the BA bit is 1, a branch to that instruction from another location is allowed. When the BA bit is 0, a branch to that instruction is not allowed.

**2.39** The octal or hexadecimal representation of the instruction follows the BA bit, and may be a single or a double word instruction. For example, the octal representation of the two-word instruction at address 00053741 (Fig. 12) is

```
1 007 00 01 020006
```

The 1 represents the BA bit, 007 00 01 represents the first word, and 020006 represents the second word. The 007 represents the OP code and 00 01 represents the Operand for the first word. The second word of the instruction is represented by 020006. The field for the second word is blank when the instruction is a one-word instruction.

The hexadecimal representation of the above two-word instruction is:

1 0701 2006

**2.40** The 2-digit number following the absolute address (for both the 3A Processor and the No. 2B ESS instructions, Fig. 14) is the CSECT (control section) number. CSECT is a pseudo-operation which identifies the beginning or extension of a control section. The 6-digit octal number following the CSECT number is the location counter (offset) within the CSECT. Each time a CSECT pseudo-operation is encountered with a new name, the counter (associated with that CSECT name) is cleared. Every time a CSECT pseudo-operation is encountered with a previously established name, the location counter is set to the next unused location associated with that CSECT name.

**2.41** For a 3A Processor instruction the encoded instruction field (Fig. 14) consists of two 6-digit octal numbers or one or two 4-digit hexadecimal numbers. The first 6-digit octal number is the first word of the encoded 3A CC instruction. The first digit in the number represents the branch-allowed (BA) bit. When the bit is equal to one, a branch to that instruction from another location is allowed; whereas, a zero means a branch to that instruction is not allowed. When an instruction is a 2-word instruction, a blank space and then another six digits are printed that represent the second word. The second octal expression is the second word of the 2-word instruction. This field is blank when the instruction is a one-word instruction.

**2.42** For a 2B ESS instruction, the encoded instruction field consists of a 3-digit octal number and a 6-digit octal number (Fig. 14). The 3-digit number is the octal representation of the high eight bits of the 24-bit No. 2B ESS instruction. The 6-digit number is the octal representation of the low 16 bits of the instruction. The first digit in this 6-bit number represents the transfer-allowed (TA) bit (similar to the BA bit used with the 3A Processor instructions).

**2.43** Refer to Sections 232-164-115 and 254-340-102 for a description of the basic instruction sets for the No. 2B ESS and the 3A Processor.

#### **Branch Address Field**

**2.44** The branch address field contains the address of the location to which control is transferred by the execution of that line of instruction, or it contains the program name where a symbol is defined. The branch address field entries may be either octal or hexadecimal numbers or name entries. If the entry is an octal or hexadecimal number, the number is the transfer address in this same program listing to which control is transferred by the instruction. If the branch address field entry is a program name, the execution of the current instruction may transfer control to the named program; however, it may be that the current instruction does not transfer control to the named program. If the instruction uses a symbol defined in the named program, control is not transferred to the named program but the name is given of the other program where the symbol is defined. The entries in the branch address field are helpful, therefore, in identifying locations (addresses) in the current program or other programs to which control is transferred or in locating program listings where a symbol used in the current program is defined.

#### **Editor Line Number Field**

**2.45** The editor line number (Fig. 12) is inserted by an editor program which may be used in the assembly of a program. This number is in decimal digits in a sequential order starting with 1 and increasing until the last line of the listing. The editor line number is interspersed with the macro or EPL expansion level number when the listing contains macro or EPL expansions. There is no line number when macros or EPL are not expanded in the listing.

#### **Macro or EPL Level Number Field**

**2.46** When a macro is called, the sequence of program instructions defined by the macro is printed whenever the macro is expanded. The printing of the system macro expansion is controlled by the MACPRNT macro options ON and OFF. If ON is specified all system macros will be expanded, and if OFF is specified the system macros are not expanded (printed). The macro instructions are identified by a macro level number (Fig. 12). The level number identifies which level of nested macro generated this instruction line. The macro level number is preceded by and followed by a dash

(-XXX-); therefore, instructions generated by a macro and printed in the listing are easily identifiable, and replace the editor line number in the editor line column. The macro level number does not affect the sequence of the editor line numbers, but is interspersed with them.

**2.47** Expansions of EPL statements are identified in the same manner as a macro expansion and have the same form.

#### Page Line Number Field

**2.48** The line number (Fig. 12) contains two decimal digits which represent the line number on the program listing page. The line numbers begin with line 01 at the top of the page and may go up to the full page capacity of 50 lines.

#### Location Symbol Field

**2.49** A location symbol is a symbolic name for a location by which branch instructions may refer to the instruction or to the block of code represented by the symbolic name. For instance, a branch instruction will refer to the symbolic name and control will be transferred to the named location.

#### Symbolic Operation Code Field

**2.50** The symbolic operation (OP) code field contains the mnemonic or symbolic representation of the operation to be performed. It may be a basic instruction, an EPL operation, a macro name, or a pseudo-operation. The various operations are explained in Section 254-340-102 and Section 254-340-100.

#### Variable or Operand Field

**2.51** The next field contains the variables or operands utilized by the operation. These may be in various forms, eg, symbolic, numerical, etc. (See Section 254-340-100.)

#### Comment Field

**2.52** The comment field is not used in the assembly or execution of the program instructions, but is used only to note operations, describe, or explain the instructions. A comment is always preceded by a sharp (#) sign at the beginning of the comment. A comment may be located on a

line to the right of an instruction or it may occupy a whole line itself. The sharp sign is used as a flag to indicate that no more executable code or assembly instruction follows the sign; however, whatever is entered after the sign is to be printed on the program listing as a comment.

#### H. Symbols Reference Tables

**2.53** The table (sometimes called the Attributes and References Table) serves as an index into the program listings for the symbols (or names) that appear in that particular program listing. The table (Fig. 15) provides an easy access to the program pages and page line numbers where the symbol appears in the current listing and is referenced to an external program if the symbol is defined in the external program. A page and line number reference is also given to provide a ready reference to the page and line number where the symbol is defined. The symbols or names appear alphanumerically under the column named NAME. The information contained in the table is formatted under column headings with meaningful names. The column headings are: VALUE, T, NAME, DEF/REF, and ATTRIBUTES AND REFERENCES in the order of appearance (left to right) across the table.

**2.54** The value of a symbol is listed under the heading VALUE within the symbol reference table. The equivalence of the symbol depends on the type of symbol and may be a numerical value, a PR name in which the symbol is defined, etc.

**2.55** The type of symbol is listed under the heading T within the symbol reference list. The type (T) information is used to distinguish symbols of different types but having the same name. The meanings of the letters shown under the T heading are given in Table C.

**2.56** The symbols are listed according to an alphanumerical collating sequence, regardless of length, under the column heading NAME. Numbers are listed after letters (letter A is low while letter Z is high followed by 0 through 9). Special characters are listed first, followed by lower case letters, upper case letters, and then numbers.

**2.57** The page number and page line number at which a symbol is defined is listed under the heading DEF/REF, and the page and line

number of each occurrence of a symbol is listed under the heading ATTRIBUTES AND REFERENCES. The page and line numbers are listed in pairs, eg, 36-15 which denotes page 36 of the listing and line number 15 on that page. The page and line numbers are listed for each appearance (or use) of the symbol in the listing. If a symbol is defined in the listing, the page and line number listed under the heading DEF/REF is the location where the symbol is defined.

**2.58** If a symbol is defined in another program listing, the name of the other listing will be given under the heading VALUE. Page and line numbers where the symbol appears in the current listing will be given under the ATTRIBUTES AND REFERENCES heading. The first appearance of the symbol (defined in another listing) in the current listing will be given under the DEF/REF heading, but the definition will be given in the other program listing. Therefore, the information under the NAME, DEF/REF, and ATTRIBUTES AND REFERENCES headings provides easy access to definitions of and to each use of each symbol used in a program listing.

**2.59** The Symbols Reference Table also contains information pertaining to symbols or names assigned to tables in memory and the layout or assignment of words and bits in the memory tables. A pictorial layout of a memory table (which may be one or more memory words) is sometimes provided in the program listing. The layout of the table is defined by a system layout macro. However, the table layout is also identified in the Symbols Reference Table by table name and item names, and by table size (TBSIZ), width (W), beginning bit position or shift factor (S), and the offset of the word (N) in the table containing the item. Figure 10 illustrates a typical pictorial layout of a memory table (CART\_LEN\_ADDR) and the identification of the table layout in the Symbols Reference Table. The table name is given under the NAME heading, and a reference (for example page 5, line 02) to the table layout definition is given under the DEF/REF heading. The number of words in the table is given by TBSIZ being equated to the number of words. Items defined as part of the table are identified by the table name appearing under the VALUE heading; and by the item name appearing under the NAME heading. The layout of the item in the table is given by the W, S, and N (width, bit position, and offset number of the word in the table,

respectively) numbers given under the ATTRIBUTES AND REFERENCES heading. For example, item CART\_LEN of CART\_LEN\_ADDR is 12 bits wide (W=C in hexadecimal), the first bit is located in bit position 4 (S=4), and is located in word 00 (N=0). The CL associated with each item is the class of the item, which is a decimal number that characterizes the table of which the item is a part.

### OTHER PROGRAM LISTING FORMATS

**2.60** Formats of program listings for the 3A Processor diagnostics are different from the format described herein because most processor diagnostics are table-driven routines. These tables are made up by various diagnostic test programs and are generated by a higher level macro language. The program section of the listings containing the data tables are different in several ways from the program section of listings described in this section; therefore, refer to Section 254-340-082 for information on the processor diagnostic listings. However, there are task monitor and interpretive routines which execute the diagnostics from the tables. These routines are in the 3A Processor language, and their program listing formats are similar to the format described in this section.

### 3. GLOSSARY

**Assembly Unit**—A collection of code that is assembled or compiled as one entity. The assembly unit is the largest physical portion of a modular program hierarchy, may or may not contain functionally related units, and is identified as a PIDENT.

**ATTRIBUTES**—Special characters reserved for specific functions or values which may be used only for the defined functions or values.

**BEGIN EPL**—The first EPL statement found in an EPL program listing; used to initialize EPL compiler and specify various options (see END\_EPL).

**CSECT**—A pseudo-operation that identifies the beginning or extension of an identifiable assembly unit or control section and assigns a location name to the identity. A CSECT is relocatable as an entity.

**Coded Line**—Any statement in SWAP or EPL containing the defined LOCATION, OPERATION, VARIABLE, and COMMENT fields.

**Comment Field**—The field that contains notes to explain the instruction, operation, etc. In the listing a note is always preceded by a sharp (#) sign. Everything to the right of the sharp sign is taken as a comment.

**DATASECT**—A system macro used to declare and open a data CSECT (which is given a label or name) for data generated by system macros. A DATASECT is relocatable.

**END EPL**—The last EPL statement; used to inform the compiler that no more EPL statements are to follow (see BEGIN\_EPL).

**Entry Point**—A labeled (named) location in an assembly unit, a program unit, a CSECT, or a block of code to which control may be transferred by a calling program, ie, the block of code may be entered at the labeled location.

**Encoded Instruction**—The octal or hexadecimal number representing the binary machine instructions and equivalent to the symbolic instruction.

**Exit**—A location at which a unit terminates and passes control to another location.

**Location Field**—The field of a formatted symbolic instruction which contains a label or name by which branch instructions may refer to the named instruction.

**Macro**—An instruction in EPL that is equivalent to a specified sequence of machine instructions, and is assigned a name by which it is called. When the macro is called, the name is expanded into the precoded lines of the macro routine. An active macro generates executable code while a declarative macro generates a data structure for an active macro.

**Mnemonic**—A functionally meaningful name assigned to represent various operations, data, locations, etc, or a combination of letters which convey the essence of an instruction (A for add, COM for complement, etc.).

**Module**—A small block of instructions within a CSECT which performs an identifiable subfunction.

**NAME**—A mandatory pseudo-operation which assigns a unique group of symbols to a PIDENT.

**Op Code**—The portion of an instruction which designates the operation that is to be performed (normally represented numerically). The op code specifies the operation that is to be performed upon the operand.

**Operand**—A value, label, or parameter to be processed by the associated op code.

**Operand Field**—The field of the instruction that contains parameters related to the operation and may contain predefined functions, symbolic variables, etc.

**Operation**—A code which directs the action of the EOS on other software, usually in a specific field of an instruction, and may be an operation code, pseudo-operation, macro, or other EPL constructions.

**PIDENT NAME**—A program identifier, which is a group of from 1 to 6 characters used to identify an assembly unit. A PIDENT number and title are normally associated with the PIDENT NAME.

**PR Number**—An identification number assigned to a PIDENT or to a collection of PIDENTS.

**Program Listing**—A printout of code assembled as a unit by the assembler program. This unit or program listing (PR) may also be referred to as a PIDENT.

**Program Unit**—A collection of code within an assembly unit which performs a well-defined function.

**Pseudo-Operation**—An operation which may be used by the assembler to control assembler activities but does not result in executable machine code.

**References**—The identification in the Attributes and References Table of the specific page number and line number where a symbol, label, name, or special character is used within an assembly unit or program unit. The page number is listed first and the line number follows. A dash (-) separates the page and line numbers.

**Prologue**—A narrative or tabular description of the assembly unit or program unit which is located at the beginning of the unit. The prologue tells the purpose of the unit and other general information about the unit.

SECTION 254-340-104

**Subroutine**—A sequence of instructions called from within another section of instructions to perform a specific function. Subroutines may be common to several programs.

**Symbol**—A series of predefined characters or mnemonics which represent binary values.

```
# THE TASK OF PROGRAM OPDATA IS TO UPDATE THE BACKUP FILES CONTAINING THE TRANSLATIONS FOR
# THE NO. 3 ESS OFFICE.

# EACH NO. 3 ESS OFFICE IS EQUIPPED WITH TWO TAPE DATA CONTROLLERS (TDCs 0 AND 1). EACH TDC
# IN THE WORKING OFFICE SHOULD HAVE A DATA CARTRIDGE INSERTED IN THE TRANSPORT AT ALL TIMES, EXCEPT
# WHEN MAINTENANCE FUNCTIONS REQUIRE ITS REMOVAL. THE TWO CARTRIDGES ARE IDENTICAL, AND EACH ONE
# CONTAINS SEVERAL FILES USED FOR VARIOUS PURPOSES. THE MOST CRITICAL FILES ARE THOSE WHICH
# ARE USED TO RELOAD MAIN STORE IN CASE A BOOTSTRAP IS DEEMED NECESSARY IN ORDER TO RETURN THE
# SYSTEM TO AN OPERATIONAL STATE. FOUR OF THESE FILES (CHECKSUM, TRNSLN, BACKDT1, AND BACKDT2) ARE
# UPDATED BY THIS PROGRAM. THE PROGRAM CONSISTS OF TWO SEPARATE PARTS: (1) THE UPDATING OF THE
# MOST RECENT COPY OF TRANSLATIONS CONTAINED IN FILE TRNSL, AND (2) THE UPDATING OF THE BACKDATED
# TRANSLATIONS IN FILE BACKDT1 OR BACKDT2.

# IN MOST CASES IN WHICH A MEMORY RELOAD IS REQUIRED, RELOADING THE MOST RECENT COPY OF
# TRANSLATIONS SHOULD BE SUFFICIENT TO RESTORE THE SYSTEM. OPDATA IS INVOKED AUTOMATICALLY EVERY 24
# HOURS TO UPDATE FILE TRNSLN AND THE ASSOCIATED DATA IN FILE CHECKSUM. IT CAN ALSO BE INVOKED
# MANUALLY WITH A TTY INPUT MESSAGE. IT IS POSSIBLE THAT THE UPDATING PROCESS COULD ABORT DUE TO
# AN I/O ERROR OR A REQUEST FROM THE MULTISCAN FUNCTION CONTROLLER (MSFC) TO ABORT. THIS COULD
# LEAVE THE OFFICE WITH A PARTIALLY UPDATED TRNSLN FILE. A SUBSEQUENT MEMORY RELOAD WOULD LOAD THE
# INVALID DATA INTO MAIN STORE. TWO PRECAUTIONS ARE TAKEN TO PREVENT THIS FROM OCCURRING: (1) THE
# UPDATE PROCESS TAKES PLACE ON ONE TAPE AT A TIME. THEN, IF AN ABORT SHOULD OCCUR AND A MEMORY
# RELOAD BECOMES NECESSARY BEFORE THE UPDATE CAN BE REPEATED SUCCESSFULLY, AT LEAST ONE TAPE WILL
# CONTAIN GOOD TRANSLATION DATA. (2) BEFORE THE UPDATE OF TRNSLN BEGINS, BIT 15 IN THE FIRST DATA
# WORD OF FILE CHECKSUM IS SET. THEN, IF THE UPDATE FAILS TO COMPLETE SUCCESSFULLY AND A MEMORY
# RELOAD BECOMES NECESSARY BEFORE THE PROBLEM IS CORRECTED, THE BIT WILL INDICATE TO THE SYSTEM
# INITIALIZATION PROGRAM THAT THE OTHER TAPE SHOULD BE USED FOR THE MEMORY RELOAD.

# IT IS POSSIBLE THAT AN ERROR IN THE UP-TO-DATE TRANSLATION DATA COULD CAUSE A MEMORY RELOAD TO
# BECOME NECESSARY. IN THIS CASE, RELOADING THE CURRENT TRANSLATION DATA WOULD NOT CORRECT THE
# PROBLEM. IT WOULD BE NECESSARY TO RELOAD MEMORY WITH AN OLDER VERSION OF TRANSLATIONS AND DO
# RECENT CHANGES TO BRING IT UP TO THE CORRECT, UP-TO-DATE STATUS. THESE OLDER VERSIONS OF
# TRANSLATIONS ARE KEPT IN FILES BACKDT1 AND BACKDT2. OPDATA IS INVOKED MANUALLY BY TTY MESSAGE
# TO DO THE UPDATES ON THESE FILES. NO CHECKSUMS ARE ASSOCIATED WITH THE DATA ON THESE FILES, AND
# THE UPDATE IS DONE ON BOTH TAPES AT ONCE. AGAIN, IF THE UPDATE WERE INTERRUPTED, A RELOAD WITH
# BACKDATED TRANSLATIONS COULD RESULT IN THE LOADING OF ERRONEOUS DATA INTO MAIN STORE. TO AVOID
# THIS PROBLEM, OPDATA ZEROES ONE OF TWO SPECIAL WORDS IN FILE BACKDT1 (CONTAINING THE DATE OF LAST
# UPDATE OF BACKDT1 AND BACKDT2, RESPECTIVELY). IF THE FILES SHOULD BE ACCESSED DURING A MEMORY
# RELOAD, A ZEROED DATE WORD INDICATES TO THE SYSTEM INITIALIZATION PROGRAMS THAT THE CORRESPONDING
# FILE SHOULD NOT BE USED FOR THE RELOAD.
```

Fig. 1a—Assembly Unit Prologue Example

```
NAME      MAICCI
# THE FUNCTION OF THIS PROGRAM IS TO INITIALIZE THE EOS AND ASSOCIATED
# TASKS.  IT MANY CASES, IT DOES NOT PERFORM THE ACTUAL INITIALIZATION,
# BUT RATHER, CALLS THE APPROPRIATE INITIALIZATION CODE AS SUBROUTINES.

# IF THE SYSTEM IS STARTED EITHER VIA A RAID ST OR CONT COMMAND, THE
# SYSTEM IS PLACED IN THE DEBUG MODE.  THE DIFFERENCES BETWEEN DEBUG
# AND NORMAL MODE ARE SLIGHT.  THEY ARE AS FOLLOWS:
#   (1) THE INITIALIZATION DB COUNTING DISPLAY ONLY OCCURS IN THE DEBUG
#       MODE.
#   (2) ON A START COMMAND, CINIT DOES NOT RETURN TO CIPL FOR THE
#       FINAL PHASE OF THE BOOTSTRAP.  IT IS ASSUMED THAT THE SYSTEM HAS BEEN
#       LOADED FROM A STR.  CARE MUST BE TAKEN SO THAT THE SYSTEM NEVER GETS
#       INTO THIS MODE IN A WORKING OFFICE.

# TWO WORDS ARE DEFINED IN MAIDTA WHICH INDICATE IF THE SYSTEM IS IN
# THE DEBUG MODE.  THESE ARE SET TO ARBITRARY VALUES TO INDICATE
# THE DEBUG MODE AS FOLLOWS:
#   MAIDDEBUG ----- SET TO DEBUG IF SYSTEM WAS STARTED WITH A ST OR
#                       CONT COMMAND.
#   MAIDSTART ----- SET TO START IF SYSTEM WAS STARTED WITH A ST COMMAND.
```

Fig. 1b—Assembly Unit Prologue Example



PROLOGUE →

```

770 01 # DESCRIPTION:
771 02 # CALLED AT THE CONCLUSION OF DIAGNOSTICS TO CHECK
772 03 # IF CU RESTORAL WAS THE CAUSE FOR RUNNING DIAGNOSTICS
773 04 # IF IT WAS AND IF DIAGNOSTICS ATP, THE OFF-LINE CU
774 05 # IS MARKED AS RESTORED.
775 06 #
776 07 # ENTRY POINTS:
777 08 # RSTCUP--IF DIAGNOSTICS ATP
778 09 # RSTCUF--IF DIAGNOSTICS FAIL
779 10 # RSTCUPF--CF=1 IF ATP, CF=0 IF FAIL
780 11 #
781 12 # ENTRY CONDITIONS:
782 13 # SEE ENTRY POINTS
783 14 #
784 15 # EXIT CONDITIONS:
785 16 # NONE
    
```

```

00053761          787 20 RSTCUPF
00053761 1 130 004          0053765 788 21 BC RSTCUP
00053762          789 22 RSTCUF
00053762          790 23 BEGIN ( )
00053762          791 24 CALL RMV_CC
00053762 1 076 01 0054102 0054102 -001- 25 BSA RMV_CC
00053764          792 26 RETURN
00053764 1 135 00 00          -001- 27 BTSA

00053765          794 29 RSTCUP
00053765          795 30 BEGIN
00053765 1 163 01 00          -002- 31 HA
00053766 063 04 0021025 CTSD 796 32 LAL RST_IP,SYSTATE,RA1
00053770          797 33 IF RST_IP THEN RGBEGIN
00053770 050 04 15          -004- 34 TBN RST_IP,S(RST_IP)
00053771 132 015          0054006 -002- 35 BNC IFS472
00053772 062 00 0021020 CTSD 798 36 LAL RO,IM_IMAGE,RAO
00053774 144 05 00          799 37 ZBS O(RAO),S(ERRI)

800 38 # THIS SUBROUTINE IS ONLY CALLED AT BASE LEVEL.
801 39 # THE EXTERNAL SOURCE THAT CAN BECOME STUCK AND CAUSE
802 40 # THE ERROR INTERRUPT TO BE REMOVED IS THE OTHER MAS
803 41 # ERRORS. SINCE CU DIAGNOSTICS PASSED, THE OTHER MAS
804 42 # IS OK. IN PARTICULAR, NONE OF ITS ERROR SIGNALS IS STUCK.
805 43 # HENCE, THE ERROR INTERRUPT CAN BE DETERMINISTICALLY ENABLED
806 44 # AT THIS POINT.

00053775 100 00 00          807 45 L RO,O(RAO)
00053776 017 14 0          808 46 LSR IM,RO
00053777 006 01 00          809 47 ZR R1
00054000 005 01 01 000026 810 48 STM R1,N(UCOROFL)(RA1),ES(UCOROFL,UCORDSR,COROFL) # CLEAR MAS
                                ERROR FLAGS
00054002 007 00 01 002020 811 50 LI RO,ES(MAS_OOS,OSA_FALT)
    
```

COMMENTS →

Fig. 2a—Program Unit Prologue Example

```

02 00304
02 00304 1 73 1 0
02 00305 CC 4 0
02 00306 CC 5 1
02 00307
02 00307 CC 1 4
02 00308 OC E C
02 00309 OC F D
02 0030A OC D 5
02 0030B 07 C 1 0110
02 0030D 61 2 5
02 0030E 56 04
02 0030F OC 00
02 00310 OC 00
02 00311 OC 00
02 00312 1 OC C E
02 00313 OC D F
02 00314 5D 1 0

```

PROLOGUE

```

01 #*****
02 # SUBROUTINE:SETFLG
03 #
04 #DESCRIPTION: THIS SUBROUTINE SETS THE SPECIFIED FLAG FOR THE SPECIFIED
05 #TASK. IT MAY BE CALLED ONLY FROM THE BACKGROUND LOOP WHICH COMPRISES
06 #MAINT.
07 #
08 #ENTRY CONDITIONS:
09 # ENTRY VIA A BSA
10 # R0=SYSTEM NAME OF TASK FOR WHICH EVENT FLAG IS TO BE SET
11 # R1=EVENT FLAG NUMBER
12 #
13 #EXIT CONDITIONS:
14 # REGISTERS ARE PRESERVED
15 # CF=1 -> MESSAGE SENT OK
16 # CF=0 -> NON EXISTENT PROCESS
17 PUBLIC SETFLG
18 SETFLG
19 HA
20 LR R4,R0
21 LP R5,R1
22 SET_FLAG NAME=R4,SET=R5
-002- 23 LR R1,R4
-001- 24 LR R14,R12
-001- 25 LR R15,R13
-002- 26 LR R13,R5
-001- 27 LI R12,X(110)
-001- 28 SVC 5
-001- 29 B **4
-001- 30 NOP
-001- 31 NOP
-001- 32 NOP
-001- 33 LR R12,R14
-001- 34 LR R13,R15
35 ETSAG

```

Fig. 2b—Program Unit Prologue Example

```

*****#
#
#                               SYNC_CLI                               #
#                               SYNCHRONOUS--CLEAR INTERRUPT          #
#
# WHAT:
#   THIS SUBROUTINE HANDLES UNANTICIPATED *IID* RESPONSES FROM BOTH  #
#   DEFINED AND UNDEFINED LINES                                     #
#
# HOW:
#   DEFINED *SLA* ARE INITIALIZED AND THE *IIDTA* IS SET TO IGNORE   #
#   FUTURE ACTIVE *IID* RESPONSES                                   #
#
# ENTRY:
#   S_CALL SYNC_CLI                                                 #
#
# ENTRY PARAMETERS:
#   RP2 = SUBROUTINE RETURN ADDRESS                                  #
#   RP4 = ADDRESS OF *SLCBP* IF LINE IS DEFINED                     #
#   RP6 = ADDRESS OF *SLCBR* IF LINE IS DEFINED                     #
#   R8 = OFFSET INTO *SIIDTA* OF SUB-CHANNEL ADDRESSED BY *IID*    #
#   R9 = DEVICE ADDRESS OF *SLA* IF LINE IS DEFINED                 #
#   R12 = BIT OF ACTIVE *IID* RESPONSE                              #
#
# EXIT:
#   S_RETURN TO INSTRUCTION AFTER S_CALL                            #
#
# EXIT PARAMETERS:
#   NONE
#
# INPUT/OUTPUT:
#   TO *SLA*
#
# REGISTER USAGE:
#   R0-R1 - UNUSED
#   RP2 - SUBROUTINE RETURN ADDRESS AND EPL FREE
#   R4-R7 - SAME AS ENTRY
#   R8 - SAME AS ENTRY
#   R9 - SAME AS ENTRY
#   R10-R11 - UNUSED
#   R12 - SAME AS ENTRY
#   R13 - UNUSED
#   RA1 - EPL FREE
#
#*****#

```

Fig. 2c—Program Unit Prologue Example

```

00054152      1109 02 COUNTING
00054152      1110 03      ST   R1,0(RA1)      # UPDATE INITQCD
00054153      1111 04      IF   R1 = 0 THEN RGBEGIN # END OF INITIALIZATION INTERVAL
00054153      030 01 01      -002- 05      TZ   R1
00054154      132 022      0054176 -001- 06      BNC  IFS553
00054155      071 01 0021024 CTSO      1112 07      STL   R1,INITLVL      # ZERO LEVEL COUNT
                                1113 08 # SYSTEM IS OUT OF CRITICAL INTERVAL RIGHT HERE
                                1114 09 # IE, NEXT INITIALIZATION WILL BE LEVEL 1
00054157      076 01 0051073 CINIT      1115 10      CALL  OPPOSTMO      # TRIGGER PRINTING OF POST-MORTEM DUMP
00054161      1 076 01 0057401 BLMMA      1116 12      CALL  INITQEND
00054163      1 006 16 00      -001- 13      BSA  INITQEND
                                1117 14      ZR   R14      # ZERO BITS 19-16 OF 20-BIT CONSTANT USED BY
                                # EXCOFLMG
00054164      007 17 01 000300      1118 16      LI   R15,ES(ISC2,ISC1) # ZERO ISC BITS NOW THAT SYSTEM HAS
                                # COMPLETED THE INITIALIZATION INTERVAL
00054166      007 00 01 107161      1119 18      LI   R0,SS_RXT*E(8):ARXF
                                1120 19 # EXCOFLMG IS A SPECIAL ENTRY TO THE INIT_OCC SUBROUTINE.
                                1121 20 # WITH R0 AND R15 SET UP AS ABOVE, THE ISC BITS WILL BE
                                1122 21 # CLEARED IN ADDITION TO THE REST OF THE INITIALIZATION.
                                1123 22 # OFF-LINE CANNOT BE INITIALIZED IF IT IS CURRENTLY IN USE BY A
                                1124 23 # PROGRAM (OSA_TBLA) OR BY A PERSON (OSM_OFL).
00054170      037 02 07 003 00      1125 24      CIRM  R2,0,S8OSA(TBLA),ES(OSM_OFL,OSA_TBLA)/ES(OSA_TBLA)
00054172      132 003      0054175 -002- 02      IF   CF THEN CALL EXCOFLMG
00054173      076 01 0060715 CSYSUB      -003- 27      BNC  IFS558
                                -002- 28 IFS558
00054175      1 017 17 17      1127 29      LSR  SS_R,R15      (Comments on a block
00054176      -001- 31 IFS553      1128 30      RGEND      of code)
00054176      1129 32 OSCHKS
                                (Comments on individual
                                lines of code)
00054176      1 033 00 02      1131 38      COM  R0,R2      # PRINT MCH FAILURE MESSAGE IF MCH WENT OUT
                                # OF SERVICE SINCE LAST ENTRY
00054177      050 00 14      1132 40      TBN  R0,S(MCH_OOS) # CF=1 IF MCH IS STILL IN SERVICE
00054200      006 00 00      CINIT      1133 41      LN   R0,ERPMCH
00054201      076 01 0051637 CINIT      1134 42      CALL ERRPRTCK
00054203      1 007 00 01 000243      -001- 43      BSA  ERRPRTCK
                                1135 44      LMCH RTNMB      # GET MANUAL SWITCH
00054205      076 01 0060551 CSYSUB      -001- 45      LI   R0,RTNMB
                                -001- 46      CALL SMCH
00054207      1 006 03 00      -002- 47      BSA  SMCH
                                1136 48      ZR   R3      # INIT FOR FOLLOWING HN R3,3 IN CASE RANGE
                                # IS SKIPPED
00054210      1137 50      IF   CF THEN RGBEGIN # USE DATA ONLY IF MCH WORKS

```

Fig. 3—Block Prologue and Comments Example

CARTU 1 SYSM A022

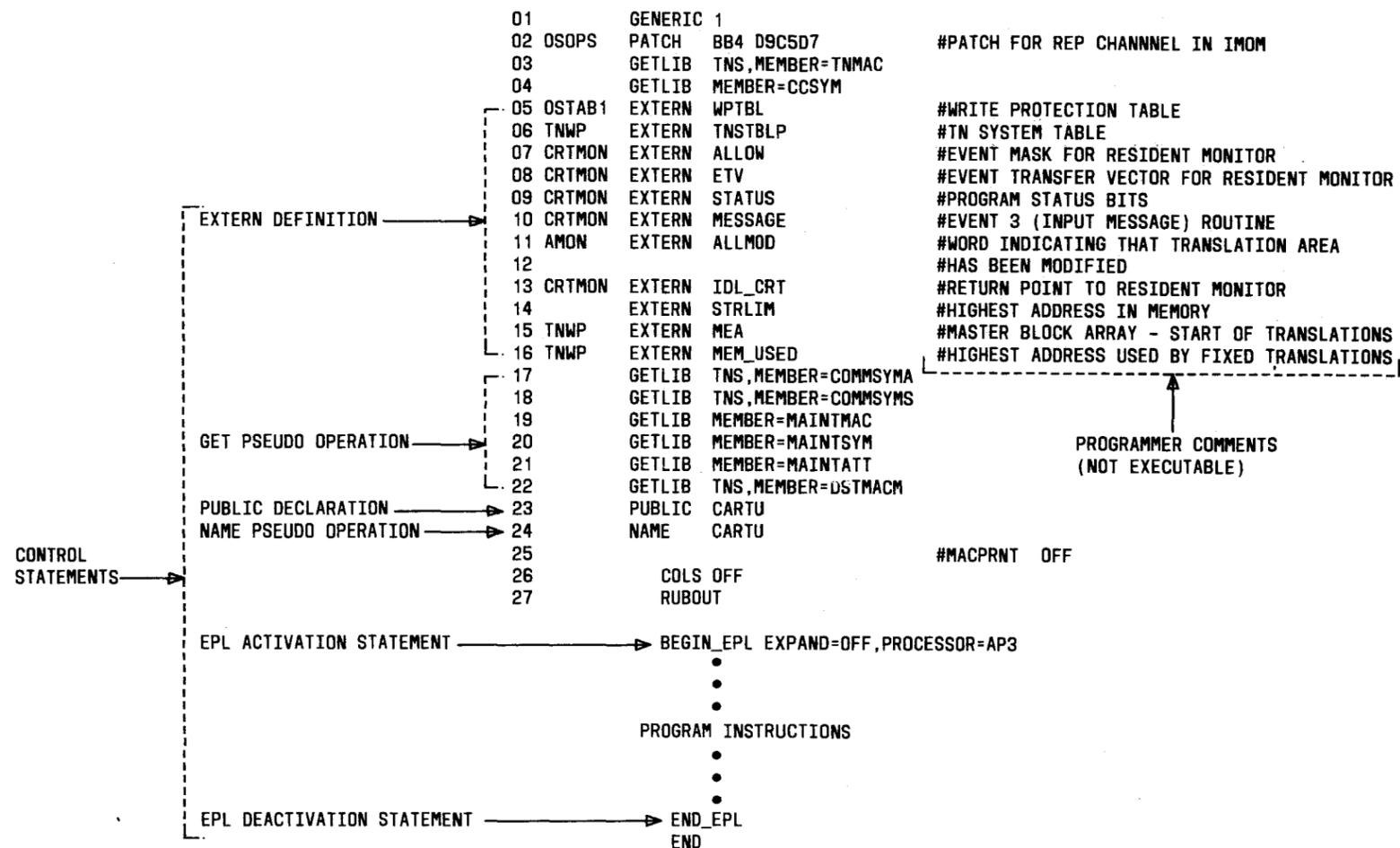


Fig. 4—Control Statements Examples

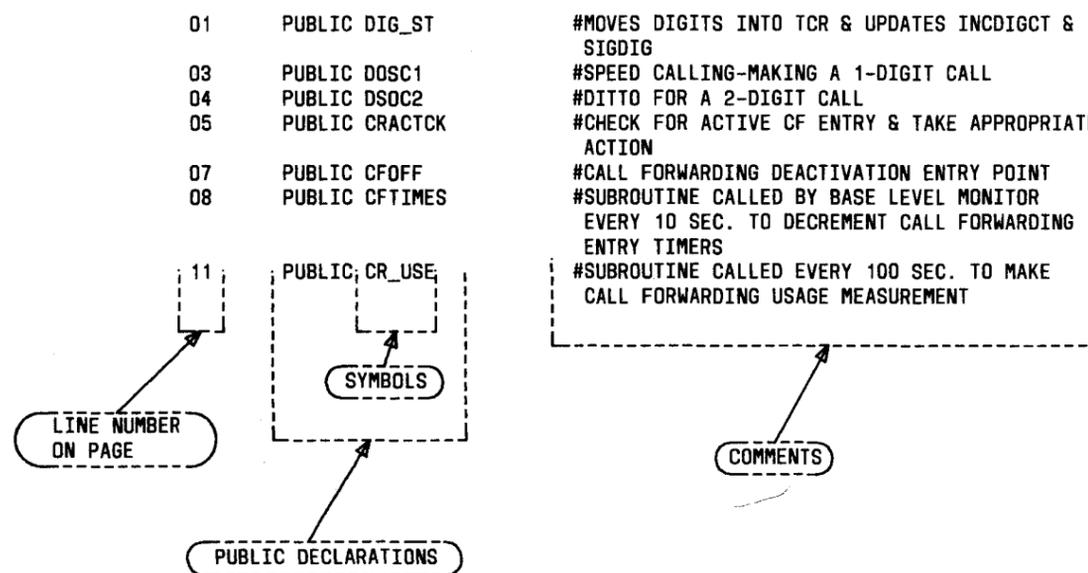


Fig. 5—PUBLIC Declaration Examples

```

01 # APPLICATION EXTERNS TQ BLMMA, INITA, TDATA, TTYAPP, AND TTYTBL APPEAR FIRST.
03     EXTERN HGAREA           # BASE OF HOLD-GET AREA
04     EXTERN TTY MFA
06 BLMMA EXTERN CLKCHGD       # SOFTWARE CLOCK IS BEING CHANGED (SUBR)
07 BLMMA EXTERN DGNCUMM       # MSF NUMBER OF CU DIAGNOSTICS (SYMBOL)
08 BLMMA EXTERN HRTBL         # DEFINITION OF TIME MONITOR HOUR FCNS (TABLE)
09 BLMMA EXTERN INITQEND      # INITIALIZATION INTERVAL JUST ENDED (SUBR)
10 BLMMA EXTERN MINTBL        # DEFINITION OF TIME MONITOR MINUTE FCNS (TABLE)
11 BLMMA EXTERN MONTBL        # DEFINITION OF MAJOR APPLICATION MONITORS (TABLE)
12 BLMMA EXTERN MSFAPCTL      # APPLICATION MSF SEQUENCE CONTROL (SUBR)
13 BLMMA EXTERN MSFMSK        # DEFINED MSFS (SYMBOL)
14 BLMMA EXTERN MSFTBL        # DEFINITION OF MSFS (TABLE)
15 BLMMA EXTERN MSFOFL        # MSFS THAT AFFECT OFF-LINE CU (SYMBOL)
16 BLMMA EXTERN PTRESET       # CONSTANT USED TO PRESET PT (SYMBOL)
17 BLMMA EXTERN SECTBL        # DEFINITION OF TIME MONITOR SECOND FCNS (TABLE)
18 BLMMA EXTERN UPD_OMAS      # MSF NUMBER OF OMAS UPDATE FCN (SYMBOL) (TABLE)
    
```

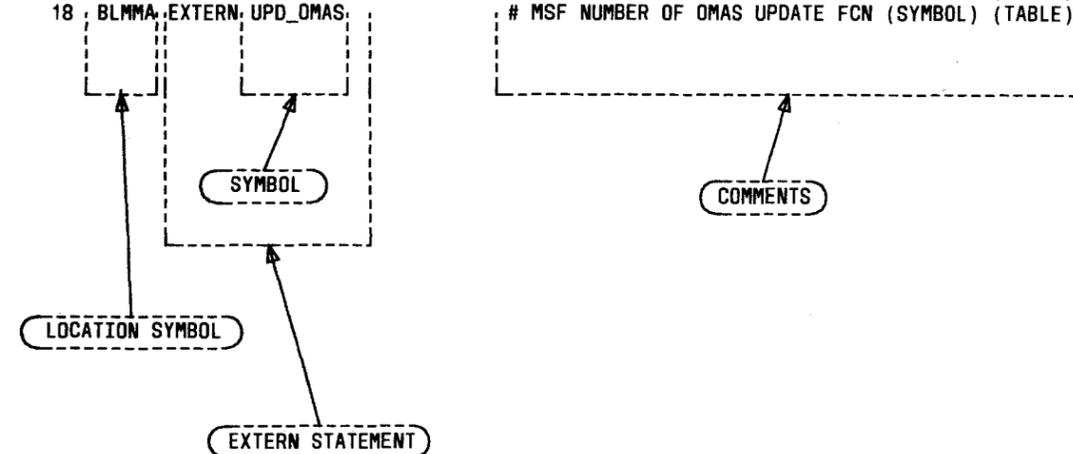
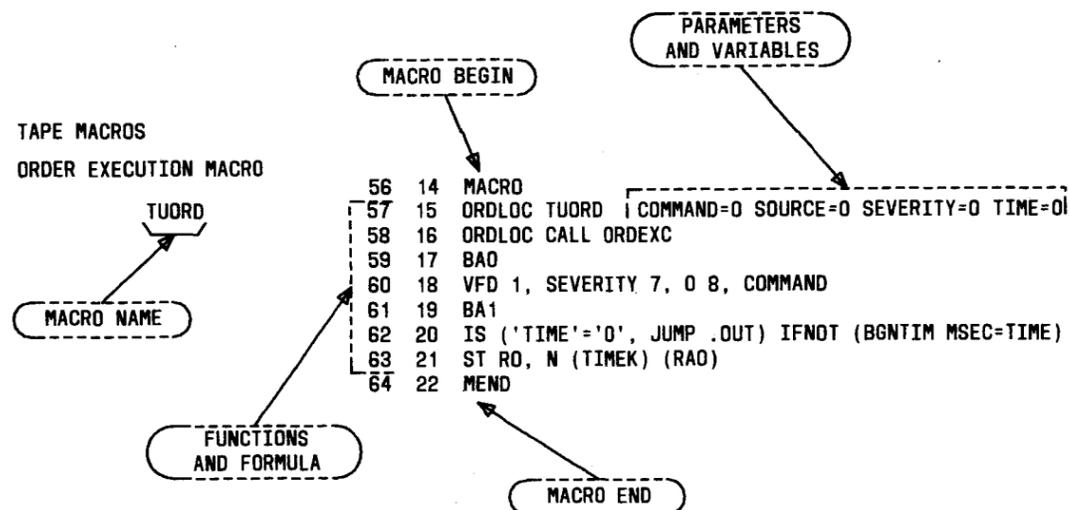


Fig. 6—EXTERN Declaration Examples



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CTAPH SYM CS31

Fig. 7—Macro Definition Example

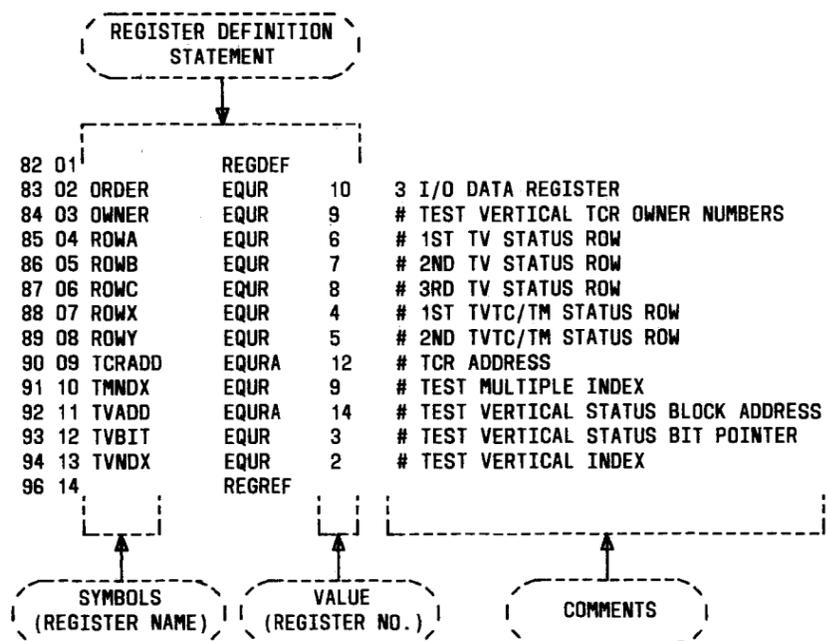


Fig. 8—Symbol Definition Example for Registers

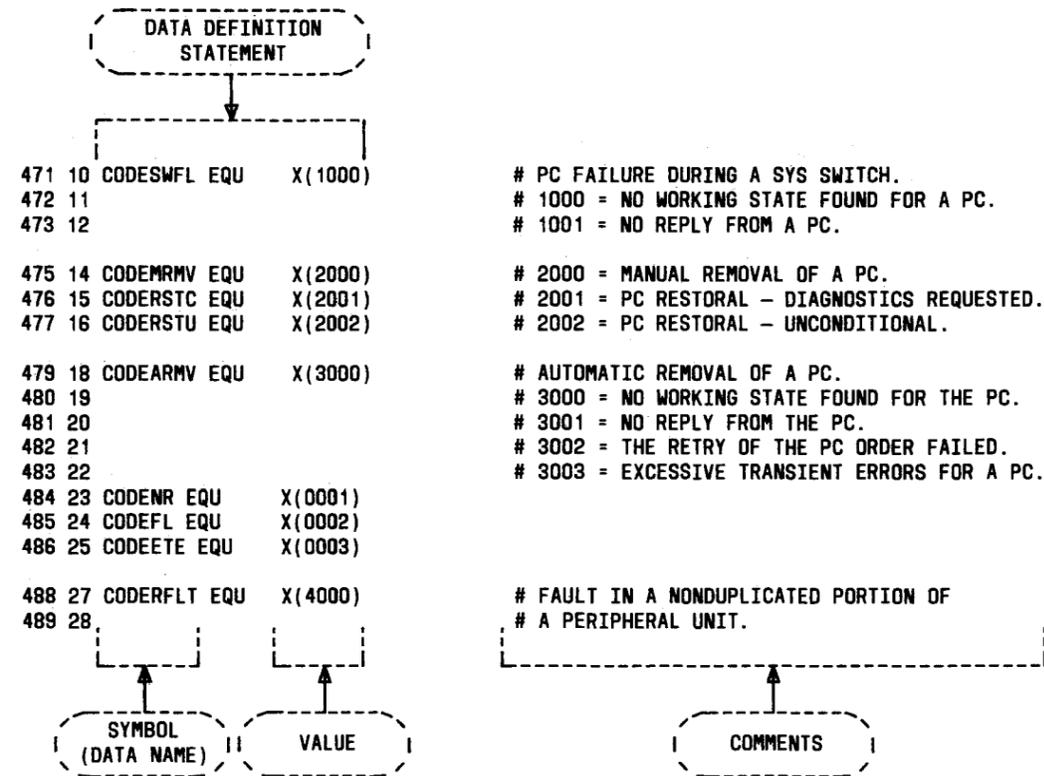
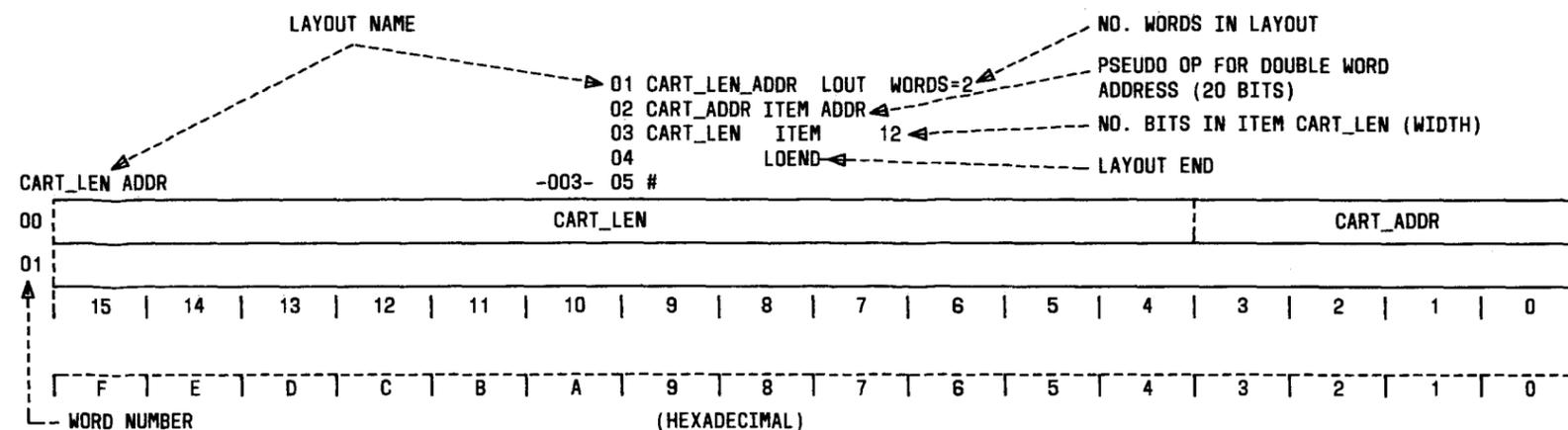


Fig. 9—Data Definition Example



VALUE	T	NAME	DEF/REF	ATTRIBUTES AND REFERENCES			
'CART_LEN_AD DR+0'	T	CART_ADDR	5-03	W=4	S=0	N=0	CL=2 19-12
'CART_LEN_AD DR+0'	T	CART_LEN	5-04	W=C	S=4	N=0	CL=2 19-11
1 00000	L	CART_LEN_ADD	5-02	TBLSIZ=2			5-35, 19-11, 19-12

ITEM NAME  
 WIDTH (NO. OF BITS) OF ITEM NAMED  
 SHIFT FACTOR (POSITION NO. OF 1ST BIT) OF ITEM NAMED  
 WORD OFFSET (NO. OF WORD CONTAINING ITEM NAMED)  
 CLASS OF ITEM NAMED  
 NAME OF TABLE AND WORD OFFSET (NUMBER) FOR ITEM NAMED

Fig. 10—Memory Table Pictorial Layout and Associated Reference in the Symbols Reference Table



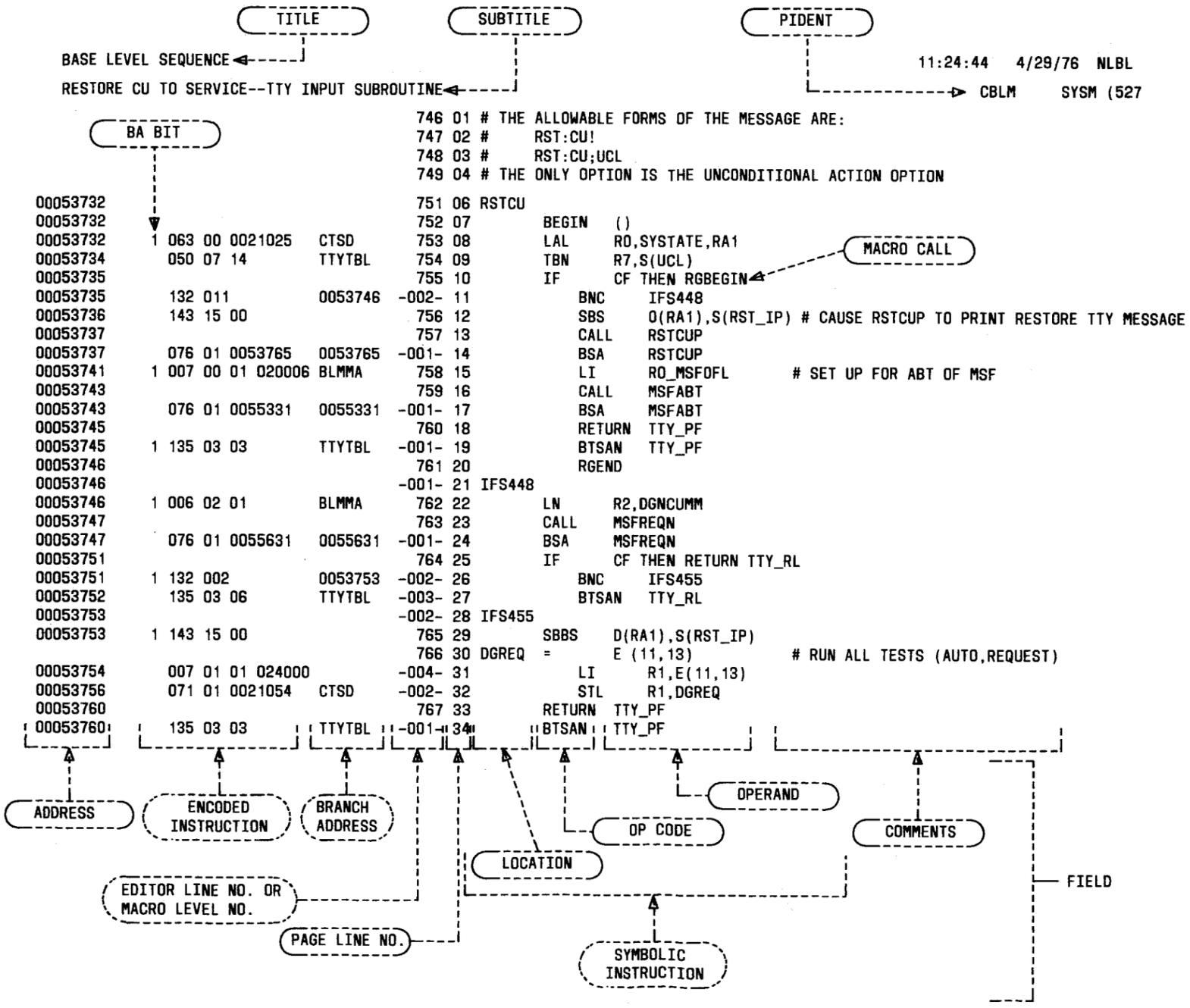


Fig. 12—Fields of Program Unit Listing—Typical Page Example

UNEXPANDED EPL LISTING

'P\_DATE', PRINT DATE 20:29:58 7/28/77 NLBL

CARTU 1 SYSM A022

02 004E4	01 P_DATE	ROUTINE	SAVE=ALL
02 004E8	02	CMCNL	OBUF,PACTION=M
02 004F0	03	DMFMT	OBUF,(WRD,WRD,WRD,WRD,DEC,ASC,ASC,ASC,ASC)
02 004F4	04	COMPOSE	<OBUF+4>,(V,F,Y)
02 004F8	05	COMPOSE	<OBUF+5>,(D,A,T,E)
	06	IF TRNS_NOW<STATUS> EQ YES_ THEN COMPOSE <OBUF+6>,(C,U,R,R,E,N,T)	
			ELSE COMPOSE <OBUF+6>,(D,L,D)
02 00505	08	COMPOSE	<OBUF+7>,(T,D,C)
02 00509	09	MULT_PRINT	OBUF,CLASS=9,FREE=R6
02 0051A	10	S_RETURN	

EXPANDED EPL LISTING

NETWORK FABRIC EXERCISE 16:24:14 10/08/76 NLBL

NFEX SYSM A213

00011445		394 39	LET <TCR+6> = R1	
00011445	1 110 01 06	-001- 40	ST R1.6(TCR)	
00011446		395 41		# EXERCISE ALL SWITCHES IN THE CONCENTRATOR GROUP.
00011446		396 43	LET <TCR+7> = (<TCR+8> = 0)	
00011446	006 02 00	-001-		# MAKE REQUEST UNCONDITIONAL.
00011447	110 02 10	-001- 45	ST R2.8(TCR)	
00011450	110 02 07	-001-		
00011451		397 47		
00011451	007 00 01 010000 TTYTBL	398 48	L1 RO.M(UCL)	
00011453		399 49	LET <TCR+9> = RO	
00011453	110 00 11	-001- 50	ST RO.9(TCR)	

Fig. 13—Unexpanded and Expanded EPL Listing—Typical Page Examples

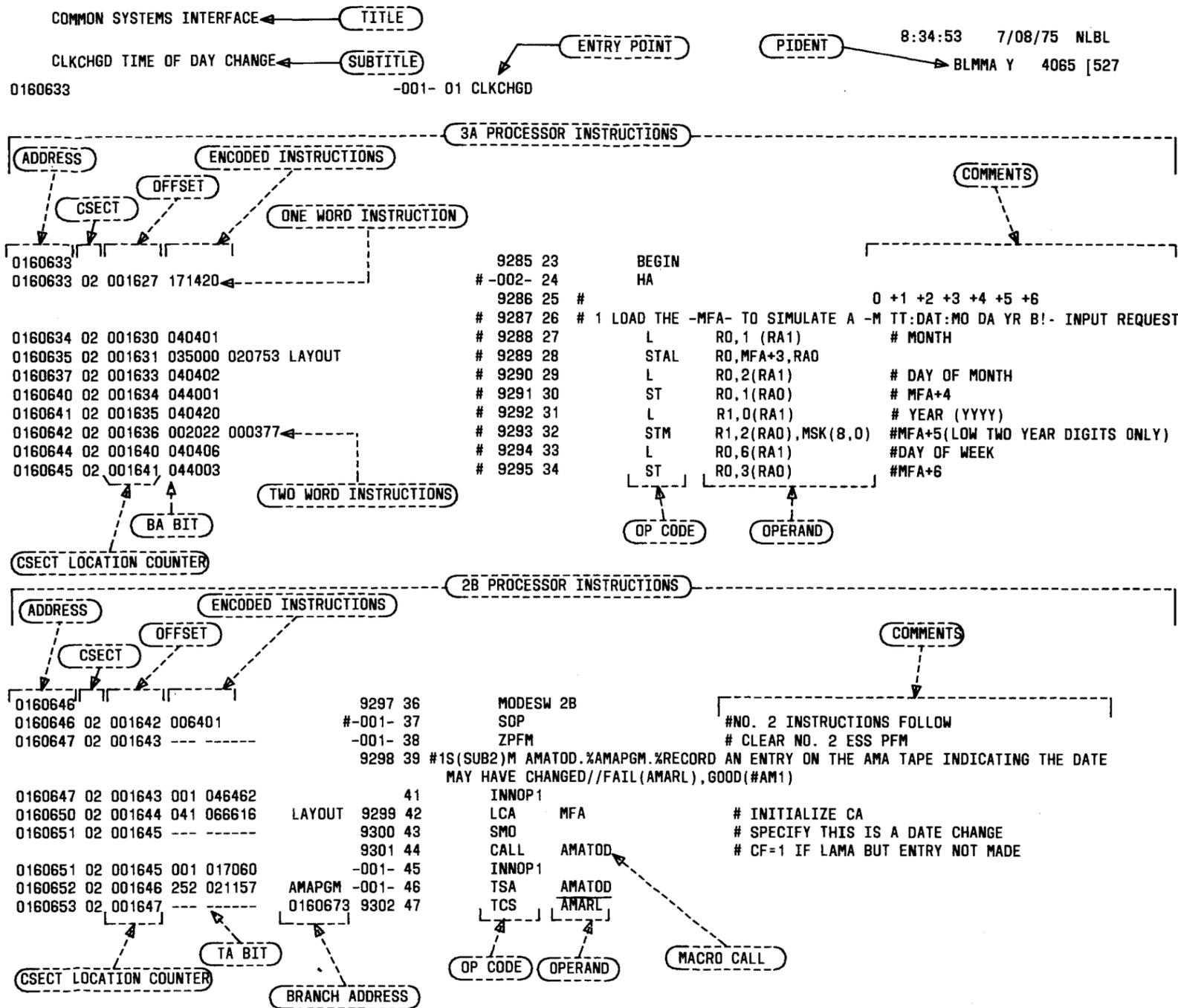


Fig. 14—Program Unit Page Listing—Typical Page Examples

-----EPILOGUE-----

14:02:05 4/10/75 NLBL

TVADM SYSM 123

VALUE	T	NAME	DEF/REF	ATTRIBUTES AND REFERENCES
				38-26, 38-42, 39-18, 39-25, 39-31, 39-32, 39-38
9	A	ROWX%SHIFT	32-14	
5	R	ROWY	9-09	28-24, 28-26, 28-28, 32-05, 38-11, 38-24
'TV_STAT+12"	T	RVGS_QC	22-02	W=2 S=2 N=C CL=2
'TV_STAT+12"	T	RVLQ_QC	22-03	W=2 S=4 N=C CL=2
0	A	RVRS	16-17	W=1 S=F N=0 CL=3
6	A	RVRSA	11-26	W=1 S=F N=6 CL=3 41-19, 41-20
8	A	RVRSB	11-32	W=1 S=F N=8 CL=3 13-23
9	A	RVRST	11-35	W=1 S=F N=9 CL=3 13-29
0	R	RO		29-11, 29-12, 29-14, 29-20, 29-41, 32-33, 32-34, 32-36, 32-39, 32-41, 32-42, 32-45, 32-46, 33-01, 35-37, 35-43, 36-04, 36-06, 36-10, 36-12, 36-13, 36-15, 36-23, 36-35, 37-06, 38-02, 38-05, 38-06, 38-38, 39-28, 41-06, 41-07, 41-08, 43-04
D	A	ROX%SHIFT	29-12	
1	R	R1		28-04, 28-05, 29-04, 29-08, 29-14, 29-19, 29-20, 29-21, 29-36, 29-37, 29-39, 32-35, 32-36, 32-37, 32-39, 32-42, 33-05, 33-06, 35-38, 35-43, 35-46, 36-03, 36-04, 36-13, 36-14, 36-17, 36-18, 36-20, 36-34, 36-45, 36-47, 37-07, 37-08, 38-29, 38-33, 38-45, 38-46, 39-08, 39-12, 39-29, 39-31, 41-03, 41-11, 41-13, 41-16, 41-23, 41-28, 43-32
4	A	R1X%SHIFT	33-05	
A	R	R10		29-15, 29-16, 29-18, 29-19, 35-11, 35-18, 35-30, 35-35, 35-48, 36-03, 36-12, 36-14, 36-15, 36-30, 36-37
4	A	R10X%SHIFT	29-18	
B	R	R11		36-22, 36-23, 36-34, 36-35, 36-36, 36-45
9	R	R9		35-10, 35-20, 36-29, 36-38, 36-48

Fig. 15—Symbols Reference Table—Typical Page Example

**TABLE A**  
**ABBREVIATIONS AND ACRONYMS**

ABBREVIATION/ACRONYM	TERM
ADDR	Address
AU	Assembly Unit
BA	Branch Allowed
CC	Central Control
CL	Class
COM	Complement
CSECT	Control Section
DATASECT	Data Section
DEF	Definition
EPL	ESS Programming Language
EQU	Equates
ESS	Electronic Switching System
LIB	Library
LR	Left Rotate
N	Index of Item
OP	Operation
PIDENT	Program Identification
PR	Program Listing
PU	Program Unit
R	Register
RA	Address Register
REF	Reference
S	Shift Factor (1st bit position of item)
T	Type of Symbol
TBLSIZ	Table Size (number of words)
W	Width of Item (number of bits)

**TABLE B**  
**NUMBER SYSTEMS**

BINARY	OCTAL	HEXADECIMAL	DECIMAL
0000	0	0	0
0001	1	1	1
0010	2	2	2
0011	3	3	3
0100	4	4	4
0101	5	5	5
0110	6	6	6
0111	7	7	7
1000	10	8	8
1001	11	9	9
1010	12	A	10
1011	13	B	11
1100	14	C	12
1101	15	D	13
1110	16	E	14
1111	17	F	15

**TABLE C**  
**TYPE SYMBOLS FROM**  
**SYMBOL REFERENCE LIST**

LETTER	TYPE OF SYMBOL
A	Absolute
B	Attribute Name
D	Program Store Data
J	Truth Valued
K	Function Name
L	Program Store Location
M	Macro
N	Pure Number
O	Machine Op
P	Address Register Pair
R	Register
S	Sequence
T	Test
U	Undefined
V	EXTERN
X	DIRECT
Y	Array Name
Z	Pseudo-operation